

G. Fountoukidou, Prof. Ioannis Pitas Aristotle University of Thessaloniki pitas@csd.auth.gr www.aiia.csd.auth.gr Version 2.2.1





DFT to FFT

- Decimation in Time (DIT) FFT
- Decimation in Frequency (DIF) FFT
- FFT Computation issues
- Goertzel Algorithm
- Bluestein Algorithm



Discrete Fourier Transform (**DFT**) of a signal x(n):

$$X(k) = \sum_{n=0}^{N-1} x(n) W_N^{nk},$$

 $x(n) = \frac{1}{N} \sum_{k=1}^{N} X(k) W_n^{-nk}.$

• *N* complex roots of unity:

$$W_N = e^{-i\frac{m}{N}}, \qquad W_N^N = 1$$

 2π

Inverse Discrete Fourier Transform (IDFT):





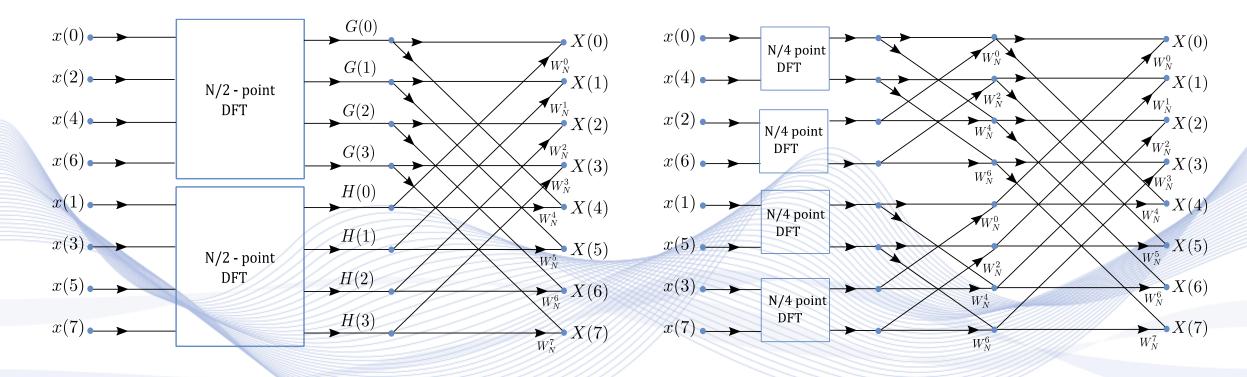
- DFT computation requires complex multiplications and additions.
- Each complex multiplication requires 4 real multiplications.
- DFT or IDFT computation by definition requires 2 for loops.
- Their computation complexity is $O(N^2)$.
- The most important advantage of DFT is that it can be calculated very fast using the *Fast Fourier Transform* (*FFT*) algorithm.





- DFT to FFT
- Decimation in Time (DIT) FFT
- Decimation in Frequency (DIF) FFT
- FFT Computation issues
- Goertzel Algorithm
- Bluestein Algorithm



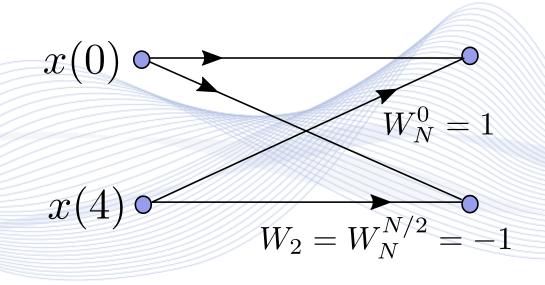


First FFT Stage for N = 8.

Artificial Intelligence & Information Analysis Lab Second FFT Stage for N = 8.



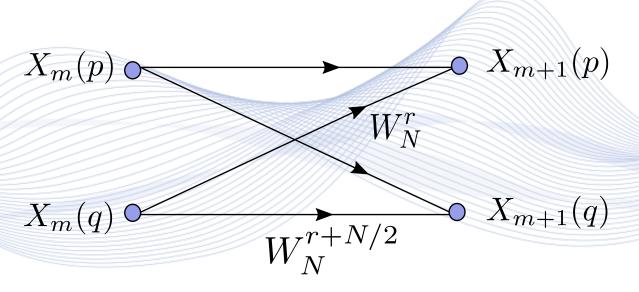
- This method results in the so-called *Decimation in Time FFT* (*DIT FFT*):
- This decimation process is continued, by breaking the calculation of the DFTs of length *N*/2 into 4 DFTs of length *N*/4, etc., until we come to DFTs of length 2:







- Each stage consists of *N*/2 basic computation structures called "*butterflies*".
- Each butterfly has 2 complex multiplication and two complex additions.







 $O(N^2)$

N = 1024

 10^{6}

 10^{4}

 $O(Nlog_2N)$

 $\blacktriangleright N$

Decimation In Time FFT

- DIT FFT has $n = \log_2 N$ steps. Each step has $\frac{N}{2}$ butterflies.
- Therefore, it requires a total of:

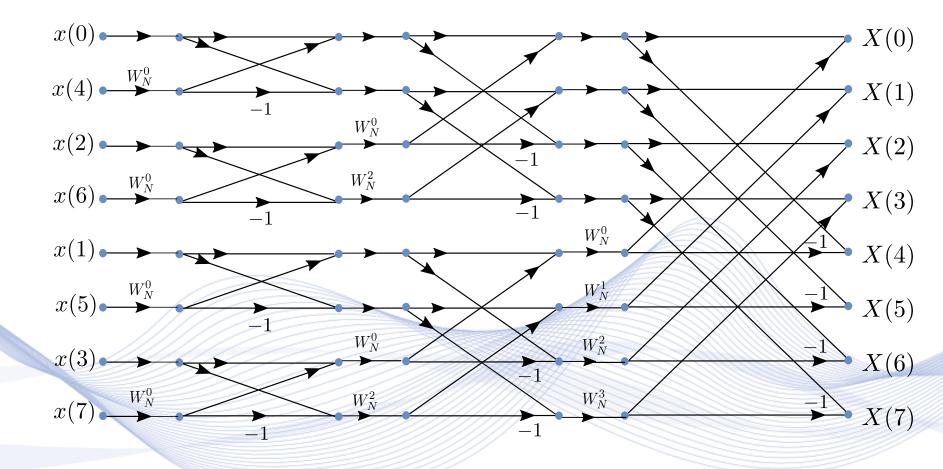
$$M = 2\left(\frac{N}{2}\right)\log_2 N = N\log_2 N.$$

complex multiplications and $Nlog_2 N$ complex additions.

• We already have a reduction in computational complexity from $O(N^2)$ to $O(N \log_2 N)$.







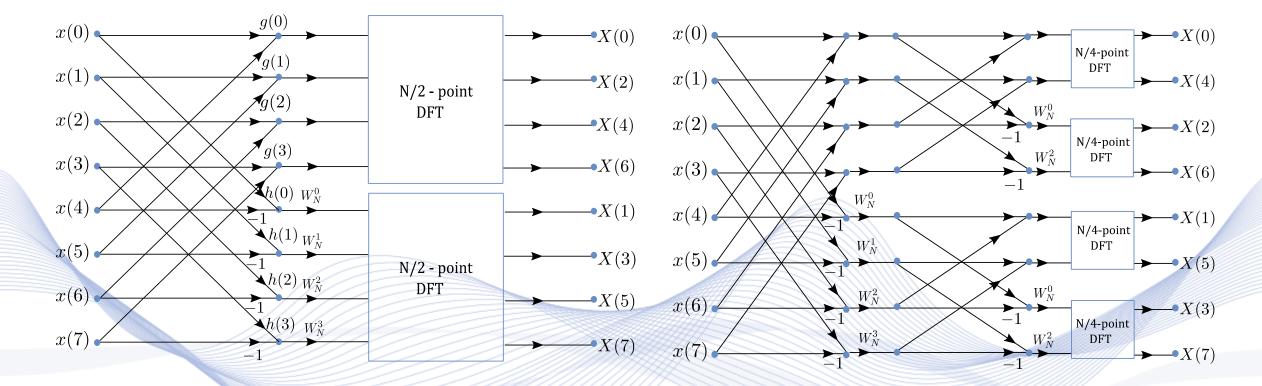
DIT FFT N = 8.





- DFT to FFT
- Decimation in Time (DIT) FFT
- Decimation in Frequency (DIF) FFT
- FFT Computation issues
- Goertzel Algorithm
- Bluestein Algorithm





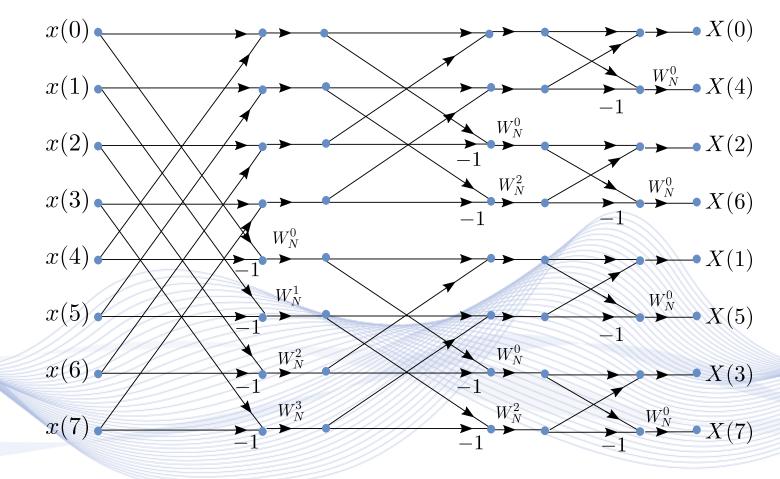
First DIF FFT stage N = 8.

Second DIF FFT stage N = 8.





Decimation in Frequency

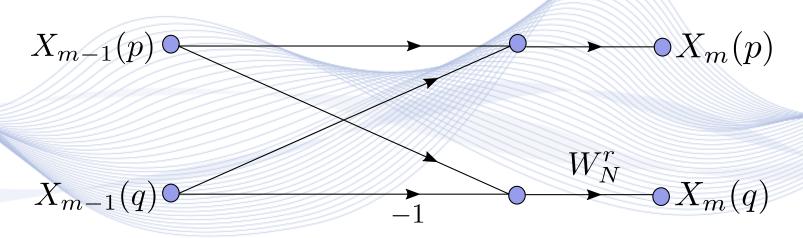


Final DIF FFT stage N = 8.



Decimation in Frequency FFT **CML**

- DIF FFT butterfly requires one complex multiplication and two complex additions.
- Therefore DIF FFT also requires $N \log_2 N$ complex additions and $(N/2) \log_2 N$ multiplications.





DIF FFT butterfly.



- DFT to FFT
- Decimation in Time (DIT) FFT
- Decimation in Frequency (DIF) FFT
- FFT Computation issues
- Goertzel Algorithm
- Bluestein Algorithm

FFT Computation Issues



There are FFT algorithms for lengths different than a power of 2:

- **Radix 4 FFT** (N is power of 4),
- Prime Factor Algorithm (PFA FFT):
 - It calculates the DFT fast, if its length is a product of prime numbers:

 $N = p_1 p_2 \dots p_n.$

• p_1, p_2, \dots, p_n : co-prime integers.





- DFT to FFT
- Decimation in Time (DIT) FFT
- Decimation in Frequency (DIF) FFT
- FFT Computation issues
- Goertzel Algorithm
- Bluestein Algorithm

Goertzel Algorithm



The IIR filter that has such an impulse response is given by:

$$H_k(z) = \frac{1}{1 - W_N^{-k} z^{-1}}.$$

Therefore X(k) calculation can be done using the IIR filter structure:

 W_N^{-k}

-1

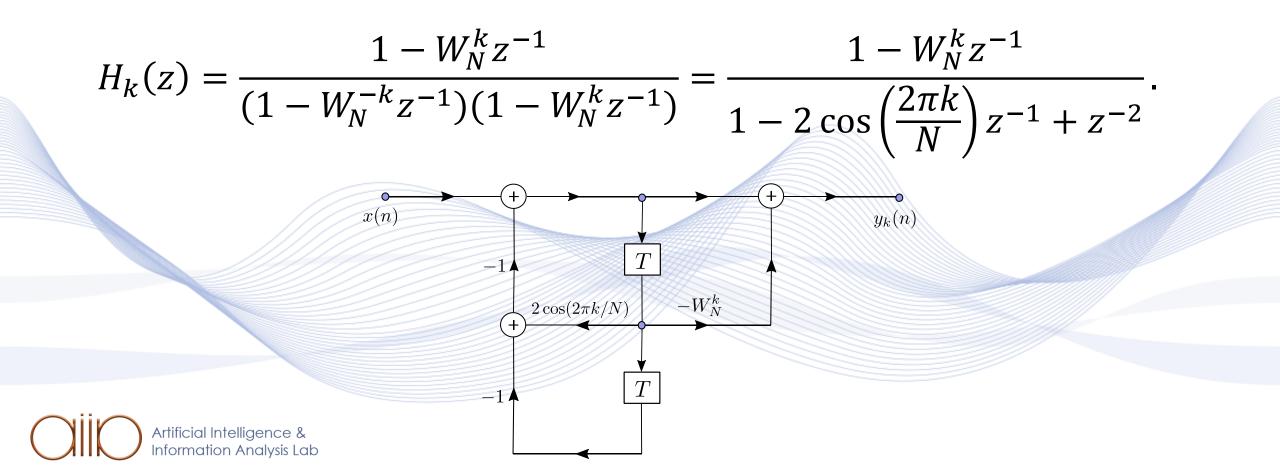
T



Goertzel Algorithm



An improved form of the *Goetzel filter* is given by the relation:



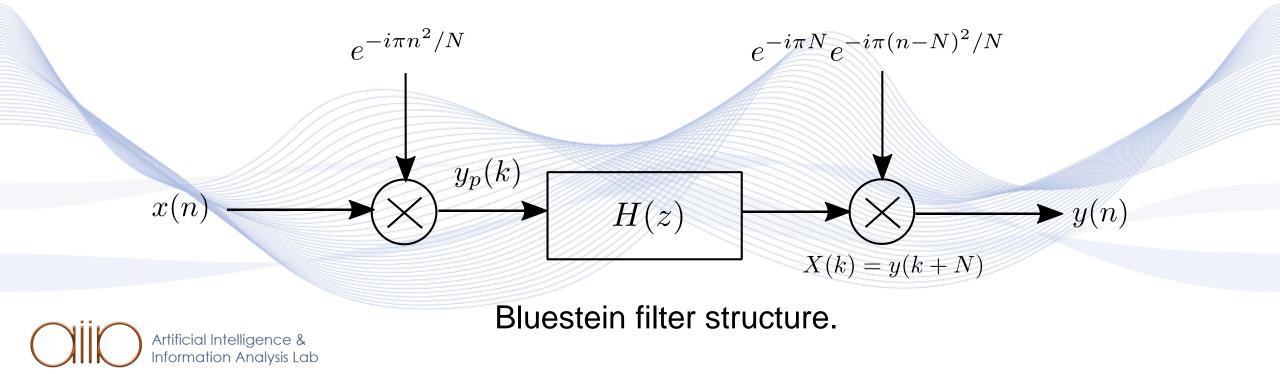


- DFT to FFT
- Decimation in Time (DIT) FFT
- Decimation in Frequency (DIF) FFT
- FFT Computation issues
- Goertzel Algorithm
- Bluestein Algorithm

Bluestein Algorithm



If we input signal $x(n)e^{-\frac{i\pi n^2}{N}}$ in the filter and multiply the output by $e^{-i\pi N}e^{-i\pi(n-N)^2/N}$, then the filter output is the DFT X(k).



Bibliography



[OPP2013] A. Oppenheim, A. Willsky, Signals and Systems, Pearson New International, 2013.

[MIT1997] S. K. Mitra, Digital Signal Processing, McGraw-Hill, 1997.

[OPP1999] A.V. Oppenheim, Discrete-time signal processing, Pearson Education India, 1999.

[HAY2007] S. Haykin, B. Van Veen, Signals and systems, John Wiley, 2007.

[LAT2005] B. P. Lathi, Linear Systems and Signals, Oxford University Press, 2005. [HWE2013] H. Hwei. Schaum's Outline of Signals and Systems, McGraw-Hill, 2013.

[MCC2003] J. McClellan, R. W. Schafer, and M. A. Yoder, Signal Processing, Pearson Education Prentice Hall, 2003.



Bibliography



[PHI2008] C. L. Phillips, J. M. Parr, and E. A. Riskin, Signals, Systems, and Transforms, Pearson Education, 2008.

[PRO2007] J.G. Proakis, D.G. Manolakis, Digital signal processing. PHI Publication, 2007.

[DUT2009] T. Dutoit and F. Marques, Applied Signal Processing. A MATLAB-Based Proof of Concept. New York, N.Y.: Springer, 2009



Bibliography



[PIT2000] I. Pitas, "Digital Image Processing Algorithms and Applications", J. Wiley, 2000.

[PIT2021] I. Pitas, "Computer vision", Createspace/Amazon, in press.

[PIT2017] I. Pitas, "Digital video processing and analysis", China Machine Press, 2017 (in Chinese).

[PIT2013] I. Pitas, "Digital Video and Television", Createspace/Amazon, 2013. [NIK2000] N. Nikolaidis and I. Pitas, "3D Image Processing Algorithms", J. Wiley, 2000.







Thank you very much for your attention!

More material in http://icarus.csd.auth.gr/cvml-web-lecture-series/

Contact: Prof. I. Pitas pitas@csd.auth.gr

